Computer Networks

Prof. Xiaoming Fu Assistants: N. Tao and D. Koll



Introduction 1-1

Course Overview

- o 25 Oct. 2012
- o 1 Nov. 2012
- o 8 Nov. 2012
- o 15 Nov. 2012
- o 22 Nov. 2012
- o 29 Nov. 2012
- o 6 Dec. 2012
- o 13 Dec. 2012
- o 20 Dec. 2012
- o 10 Jan. 2013
- o 17 Jan. 2013
- o 24 Jan. 2013
- o **31 Feb. 2013**
- o 7 Feb. 2013

Introduction & Layering Link Layer I Link Layer II Network Layer I Network Layer II; Routing I Network Layer III; Routing II; Mobility Transport Layer I Transport Layer II **Networked Multimedia** Quality of Service **Network Security I** Network Security II **Questions & Answers Session**

Written Examination



Excercises

• Contact e-mail:

koll@cs.uni-goettingen.de

- Homework exercises will be handed out regularly after class and are in the wiki.
- Students are encouraged to work on their own and solve the homework exercises to prepare for the final exam.
- Solutions will be presented one week later after class. Thursdays 12:00 – 13:00 in the lecture room.





• The grading is as follows:

100% Final exam!

All important information (click on Computer Networks)

wiki.net.informatik.uni-goettingen.de



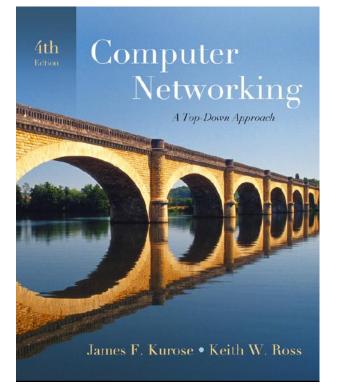
Chapter 1 Introduction

This lecture is based on the book:

Computer Networking: A Top Down Approach 4th edition. Jim Kurose, Keith Ross, Addison-Wesley, July 2007.

Alternative textbook:

- A. Tanenbaum, "Computer Networks", 5th edition, Prentice Hall, 2010
- D. Comer, "Computer Networks and Internets", 5th edition, Prentice Hall, 2008





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Chapter 1: roadmap

- 1.1 What is the Internet?
- 1.2 Network edge
 - end systems, access networks, links
- 1.3 Network core
 - circuit switching, packet switching, network structure
- 1.4 Delay, loss and throughput in packet-switched networks
- 1.5 Protocol layers, service models
- 1.6 History



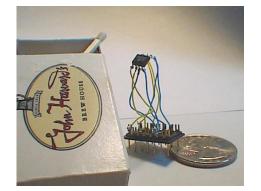
"Cool" internet appliances



IP picture frame http://www.ceiva.com/



Web-enabled toaster + weather forecaster



World's smallest web server http://www-ccs.cs.umass.edu/~shri/iPic.html

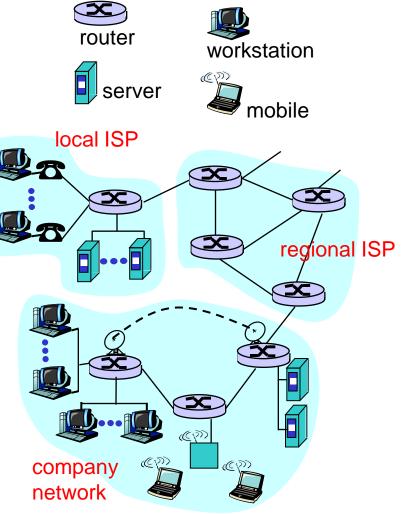


Internet phones



What's the internet? A close look...

- millions of connected computing devices: *hosts, end-systems*
 - PCs, workstations, servers
 - o PDAs, phones, toasters
 - running network apps
- o communication links
 - fiber, copper, coax, radio, satellite
 - transmission rate = bandwidth
- *routers:* forward packets (chunks of data)



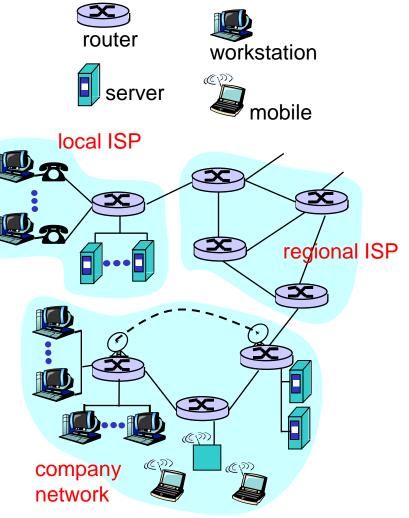
Introduction

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What's the internet? ... and closer

- protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt
- Internet: "network of networks"
 - o loosely hierarchical
 - public Internet versus private intranet
- o Internet standards
 - RFC: Request for Comments
 - IETF: Internet Engineering Task Force



Introduction

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What's a protocol?

human protocols:

- o "what's the time?"
- o "I have a question"
- o introductions
- ... specific msgs sent ... specific actions taken when msgs received, or other events

network protocols:

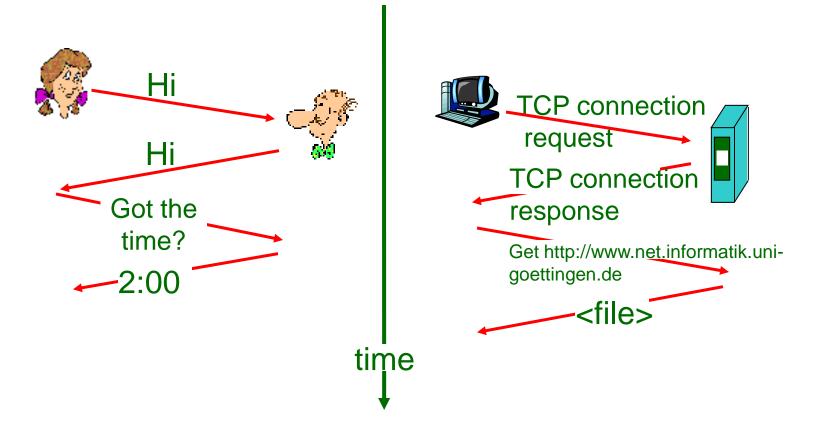
- machines rather than humans
- all communication activity in Internet governed by protocols

protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission, receipt



What's a protocol?

a human protocol and a computer network protocol:





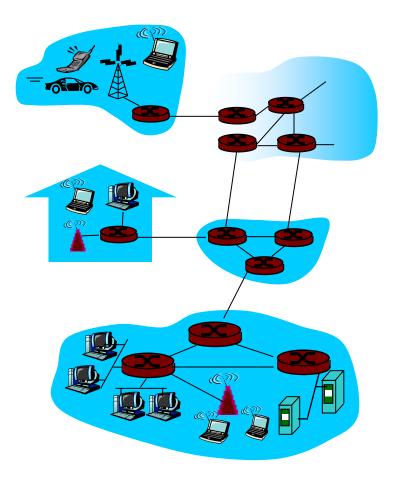
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A closer look at network structure:

- network edge:
 applications and
 hosts
- access networks, physical media: wired, wireless communication links
- o network core:
 - interconnected routers
 - network of networks





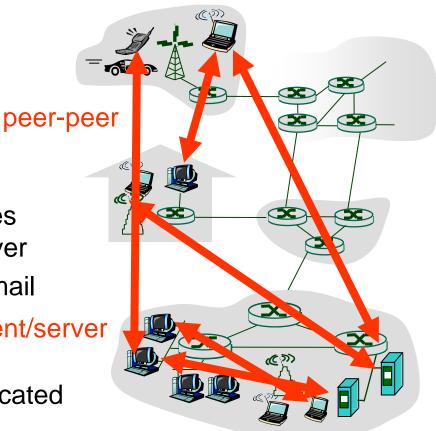
The network edge:

o end systems (hosts):

- run application programs
- e.g. web, email
- at "edge of network"

client/server model

- client host requests, receives service from always-on server
- e.g. web browser/server; email client/server
 client/server
- o peer-peer model:
 - minimal (or no) use of dedicated servers
 - e.g. Skype, BitTorrent

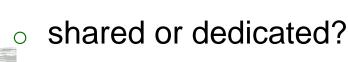


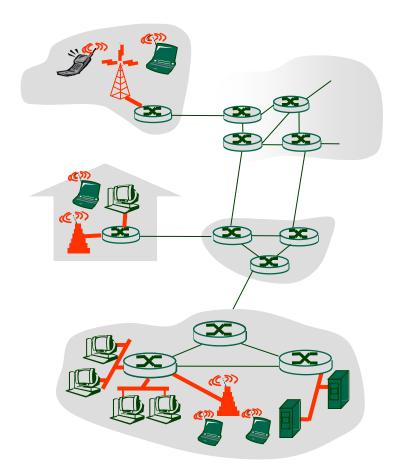
Access networks and physical media

- Q: How to connect end systems to edge router?
- residential access nets
- institutional access
 networks (school,
 company)
- mobile access networks

Keep in mind:

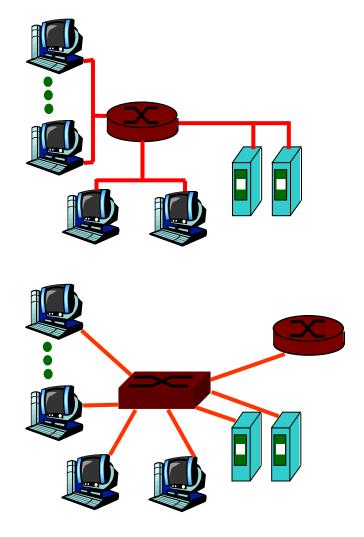
 bandwidth (bits per second) of access network?





Example: Company access: local area networks

- company/univ local area network (LAN) connects end system to edge router (example: our GöNet)
- Ethernet:
 - 10 Mbs, 100Mbps,
 1Gbps, 10Gbps Ethernet
 - modern configuration:
 end systems connect into
 Ethernet switch
- LANs: will be discussed in detail throughout this lecture

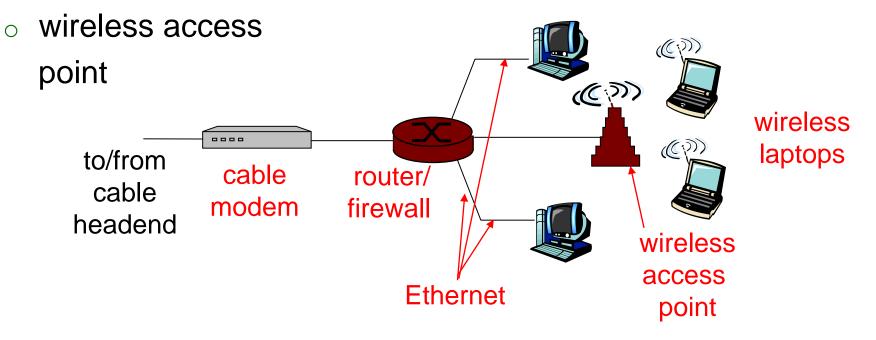




Example: Home networks

Typical home network components:

- DSL or cable modem
- router/firewall/NAT
- Ethernet





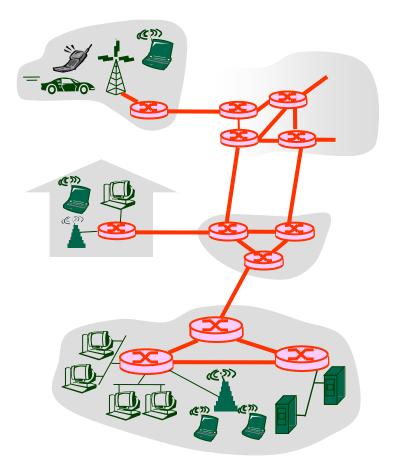
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The Network Core

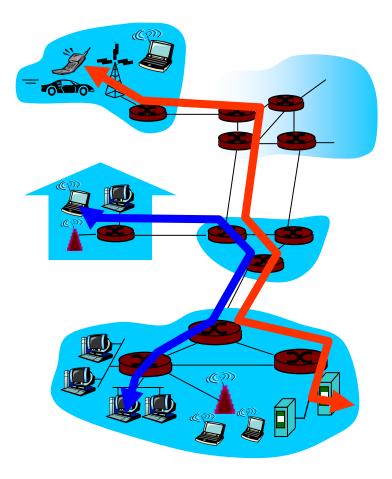
- mesh of interconnected routers
- <u>the</u> fundamental question: how is data transferred through net?
 - circuit switching: dedicated circuit per call: telephone network
 - packet-switching: data sent through a network in discrete "chunks"





Network Core: Circuit Switching

- End-end resources reserved for "call"
- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- o call setup required



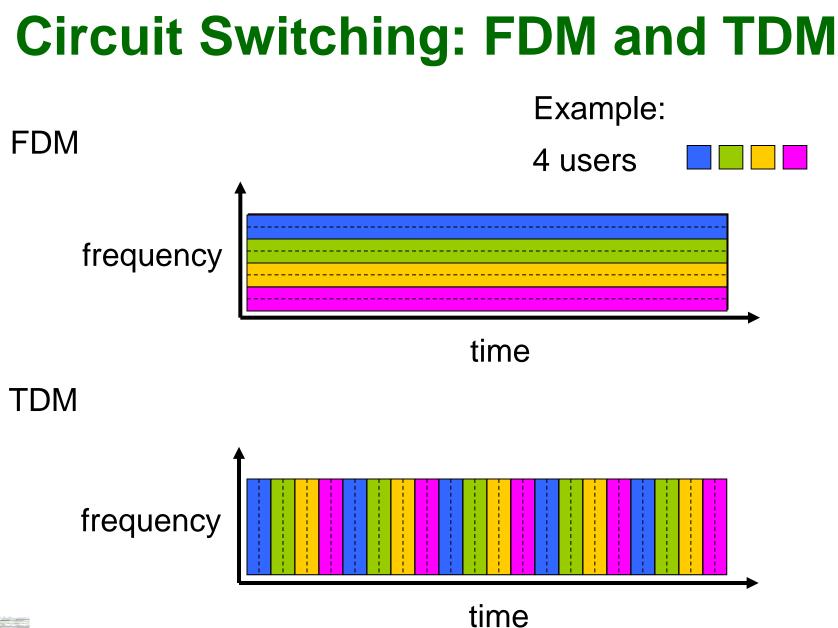


Network Core: Circuit Switching

- network resources (e.g., bandwidth) divided into "pieces"
- pieces allocated to calls
- resource piece *idle* if not used by owning call (no sharing)

- dividing link bandwidth into "pieces"
 - frequency division
 - $_{\circ}$ time division





Network Core: Packet Switching

each end-end data stream divided into *packets*

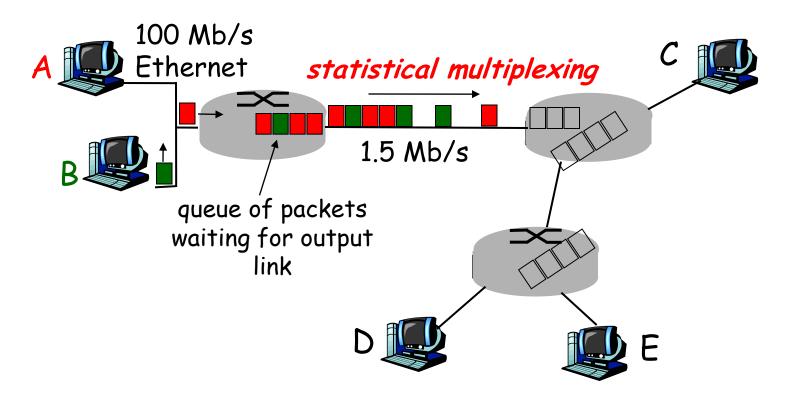
- user A, B packets share network resources
 - Sequence of sending packets does not have fixed pattern → statistical multiplexing
- each packet uses full link bandwidth
- resources used as needed

Bandwidth division into "pieces" Dedicated allocation Resource reservation resource contention:

- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward:
 packets move one hop at a time
 - Node receives complete packet before forwarding

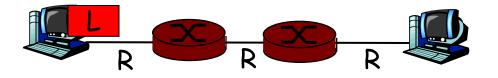


Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern, bandwidth shared on demand → statistical multiplexing.
TDM: each host gets same slot in revolving TDM frame.

Packet-switching: store-and-forward



- takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- delay = 3L/R (assuming zero propagation delay)

Example:

- \circ L = 7.5 Mbits
- R = 1.5 Mbps
- transmission delay = 15
 sec

Note:

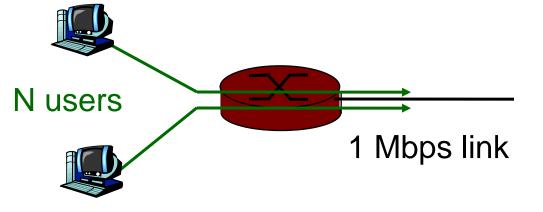
- In order to be more efficient, large packets are usually segmented into smaller packets
- → Can you explain why?



Packet switching versus circuit switching

Packet switching allows more users to use network!

- 1 Mb/s link
- each user:
 - 100 kb/s when "active"
 - active 10% of time
- o *circuit-switching:*
 - 10 users
- o packet switching:
 - with 35 users, probability
 > 10 active at same time is low



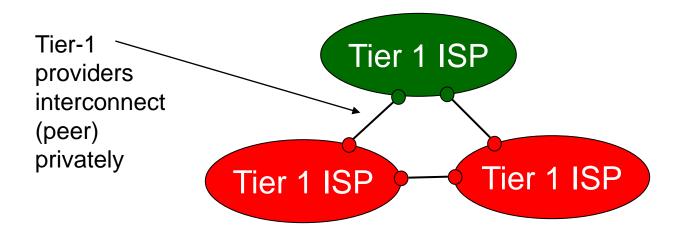
Packet switching versus circuit switching

Is packet switching better than circuit switching?

- great for bursty data
 - \circ resource sharing
 - simpler, no call setup
- excessive congestion: packet delay and loss
 - protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
 - bandwidth guarantees needed for audio/video apps
 - $_{\circ}~$ still an unsolved problem

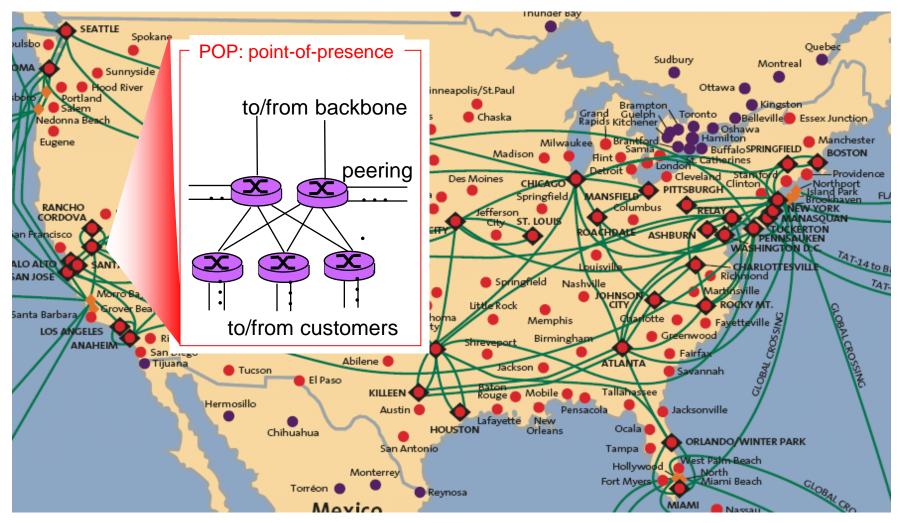


- roughly hierarchical
- at center: "tier-1" ISPs (e.g., Verizon, Sprint, AT&T, Cable and Wireless), national/international coverage
 - $_{\circ}$ treat each other as equals





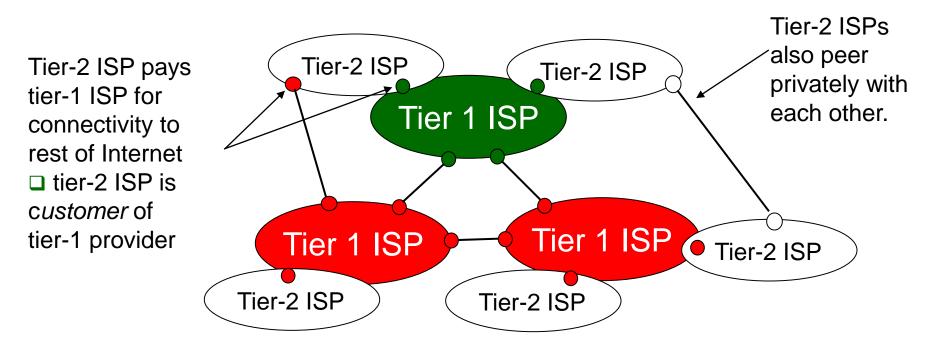
Tier-1 ISP: e.g., Sprint





• "Tier-2" ISPs: smaller (often regional) ISPs

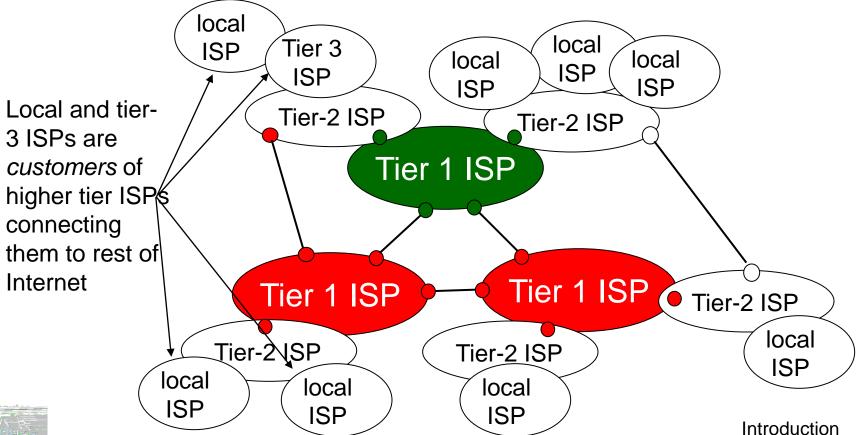
• Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



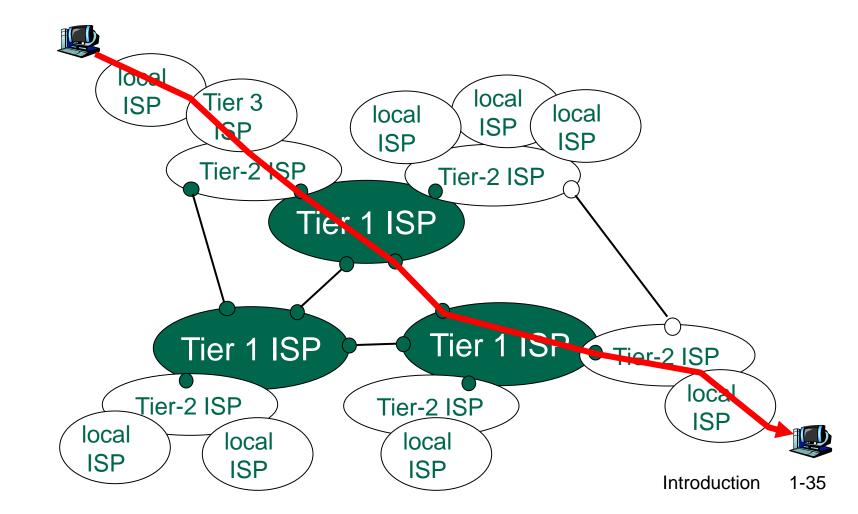


"Tier-3" ISPs and local ISPs

last hop ("access") network (closest to end systems)



• a packet passes through many networks!





Chapter 1: roadmap

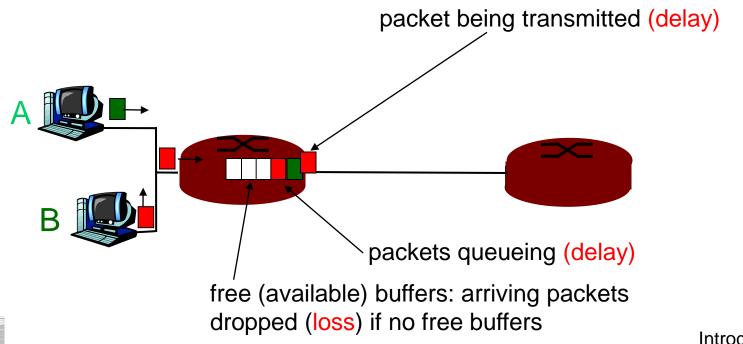
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How do loss and delay occur?

packets queue in router buffers

- packet arrival rate to link exceeds output link capacity
- o packets queue, wait for turn

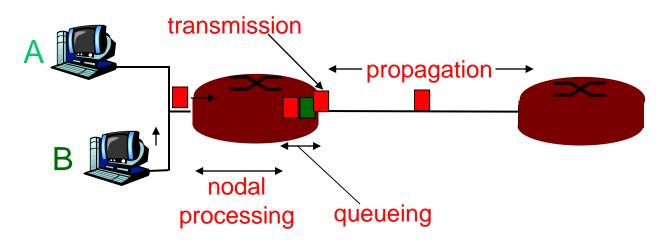


Four sources of packet delay

- 1. nodal processing:
 - check bit errors
 - determine output link

• 2. queueing

- time waiting at output link for transmission
- depends on congestion level of router



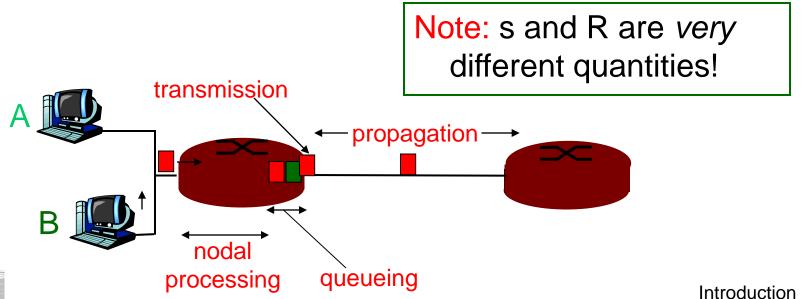


Delay in packet-switched networks

- 3. Transmission delay:
- R=link bandwidth (bps)
- L=packet length (bits)
- time to send bits into link = L/R

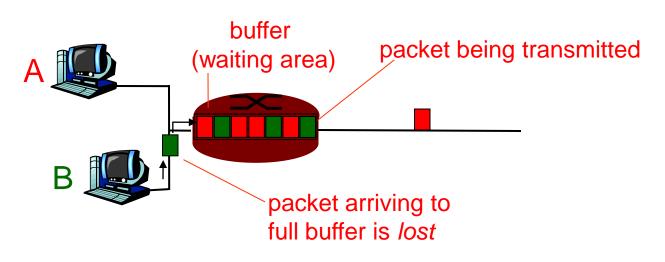
- 4. Propagation delay:
- \circ d = length of physical link
- \circ s = propagation speed in medium (~2x10⁸ m/sec)

$$\circ$$
 propagation delay = d/s



Packet loss

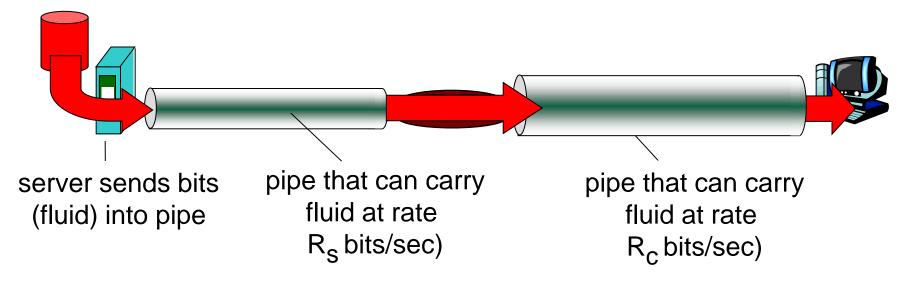
- queue (aka buffer) preceding link in buffer has finite capacity
- packet arriving to full queue dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not at all





Throughput

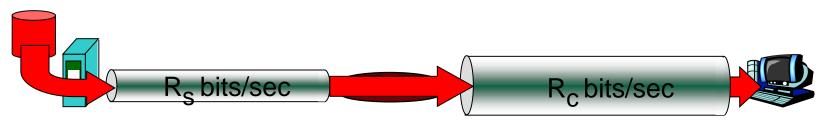
- *throughput:* rate (bits/time unit) at which bits transferred between sender/receiver
 - o *instantaneous:* rate at given point in time
 - average: rate over longer period of time





Throughput (more)

 $\circ R_s < R_c$ What is average end-end throughput?



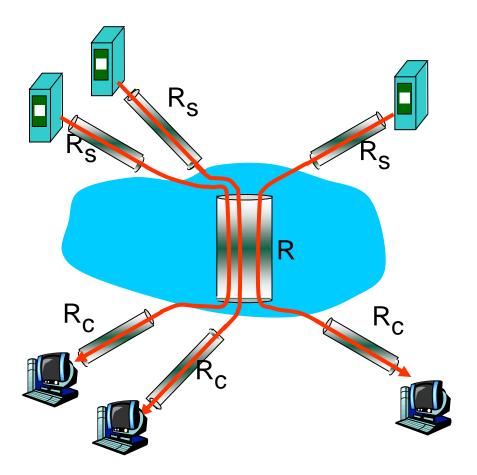
• $R_s > R_c$ What is average end-end throughput? R_s bits/sec R_c bits/sec R_c bits/sec

bottleneck link

link on end-end path that constrains end-end throughput

Throughput: Internet scenario

- per-connection endend throughput: min(R_c,R_s,R/10)
- in practice: R_c or R_s
 is often bottleneck



10 connections (fairly) share backbone bottleneck link R bits/sec



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Protocol "Layers"

Networks are complex!

- many "pieces":
 - o hosts
 - o routers
 - links of various media
 - $_{\circ}$ applications
 - o protocols
 - ₀ hardware, software

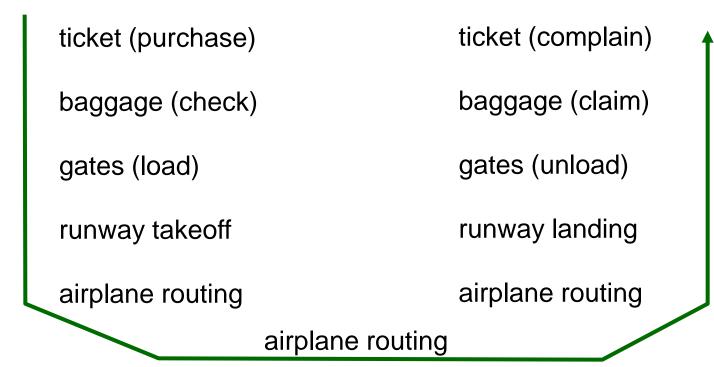
Question:

Is there any hope of organizing structure of network?

Or at least our discussion of networks?



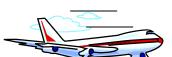
Organization of air travel



a series of steps



Layering of airline functionality



ticket (purchase)		ticket (complain)	ticket
baggage (check)		baggage (claim	baggage
actes (lead)			
gates (load)		gates (unload)	gate
runway (takeoff)		runway (land)	takeoff/landing
airplane routing	airplane routing airplane routing	airplane routing	airplane routing
departure	intermediate air-traffic	arrival	

airport

Layers: each layer implements a service

control centers

- via its own internal-layer actions
- $_{\circ}~$ relying on services provided by layer below



airport

Why layering?

Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
 - layered reference model for discussion
- modularization eases maintenance, updating of system
 - change of implementation of layer's service transparent to rest of system
 - e.g., change in gate procedure doesn't affect rest of system
- o layering considered harmful?



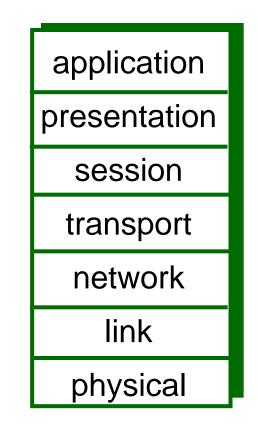
Internet protocol stack

- application: supporting network applications
 - FTP, SMTP, HTTP
- transport: process-process data transfer
 - TCP, UDP
- network: routing of datagrams from source to destination
 - IP, routing protocols
- link: data transfer between neighboring network elements
 - PPP, Ethernet
 - physical: bits "on the wire"

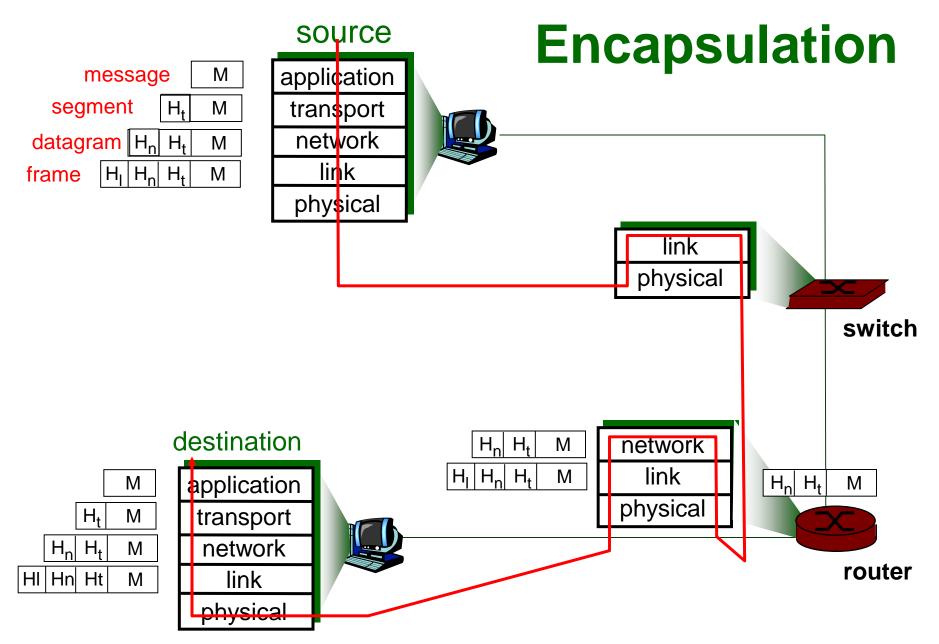
_		
	application	
	transport	
	network	
	link	
	physical	

ISO/OSI reference model

- presentation: allow applications to interpret meaning of data, e.g., encryption, compression, machinespecific conventions
- session: synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
 - these services, *if needed*, must be implemented in application
 - o needed?









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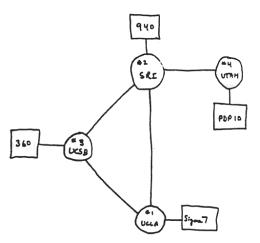


1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packet-switching
- 1964: Baran packetswitching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

o **1972**:

- ARPAnet public demonstration
- NCP (Network Control Protocol) first host-host protocol
- o first e-mail program
- ARPAnet has 15 nodes





1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- ate70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture



1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control

- new national networks:
 Csnet, BITnet, NSFnet,
 Minitel
- 100,000 hosts
 connected to
 confederation of
 networks



1990, 2000's: commercialization, the Web, new apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
 - hypertext [Bush 1945, Nelson 1960's]
 - HTML, HTTP: Berners-Lee
 - o 1994: Mosaic, later Netscape
 - late 1990's: commercialization of the Web

Late 1990's – 2000's:

- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps



Introduction: Summary

Covered a "ton" of material!

- Internet overview
 - Incl. Internet / ISP strucuture
- o what's a protocol?
- network edge, core, access network
 - packet-switching versus circuit-switching
 - Internet structure
- performance: loss, delay, throughput
- layering, service models
- o history

You (should ;) now have:

- context, overview, "feel" of networking
- more depth, detail to follow!