

Exercise 3

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Switch Learning Process

Q1: Please briefly describe the learning process that a switch uses to fill its table.

The steps taken place when a frame received:

Source Addr.	Interface #	TTL
12-34-56-78-9A-BC	1	t1

1. Add a new entry into the table. (**step 2**)
2. **if** Is the entry started with destination address of the frame found in the switch table ?
 - if** Is the interface number of the new found entry the same with the interface number the frame arrived in ?(Or is they from the same sub-network ?)
 - Drop the frame (**filtering**)
 - else**
 - Forward the frame on interface indicated (**forwarding / step 4**)
 - end**
- else**
 - flood (**step 3**)
- end**

Functionality 1: **Filtering**
Functionality 2: **Forwarding**

Step 1: to initialize the table as empty
Step 2: to build new entry
Step 3: to flood;
Step 4: to forward
Step 5: to delete entry

Hubs, Switches, Routers...

Q2: What are the main purposes of the following devices: hub, switch, and router?

What are the differences between a switch and a router?

	Hub	Switch	Router
Belong to Which layer?	1 physical layer; repeat bits	2 link layer; filter and forward frame	3 network layer; deliver datagram
Can isolate frame collision	No	Yes	Yes
Can be used to connect different Ethernet technologies	No	Yes	Yes
Multiple access protocol	No	Yes	?
Speed	1	2	3
Plug and play	Yes	Yes	No
Firewall protection	No	No	Yes
Shortest path	☹️	No	Yes

PPP Requirements

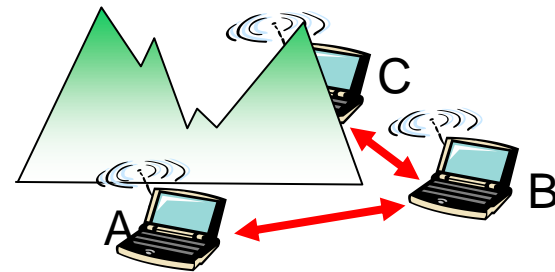
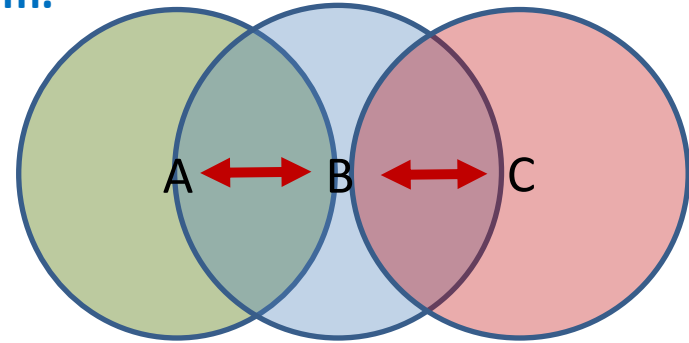
Q3: Why does PPP not need to provide error correction/recovery, flow control, etc.?

- Concept of Layering:
 - Error correction/recovery, flow control, delivery order are all delegated to the upper layers
- That means: PPP only responsible for
 - Framing of packets arriving from upper layer
 - Detection of data errors
 - Detection of link failure

Hidden Terminal Problem

Q4: Please explain the Hidden Terminal Problem.

- Appears in **wireless** networks
 - Two nodes that are not visible to each other (A,C) try to communicate to a node (e.g., an AP) visible to both (B) at the same time -> **interference**

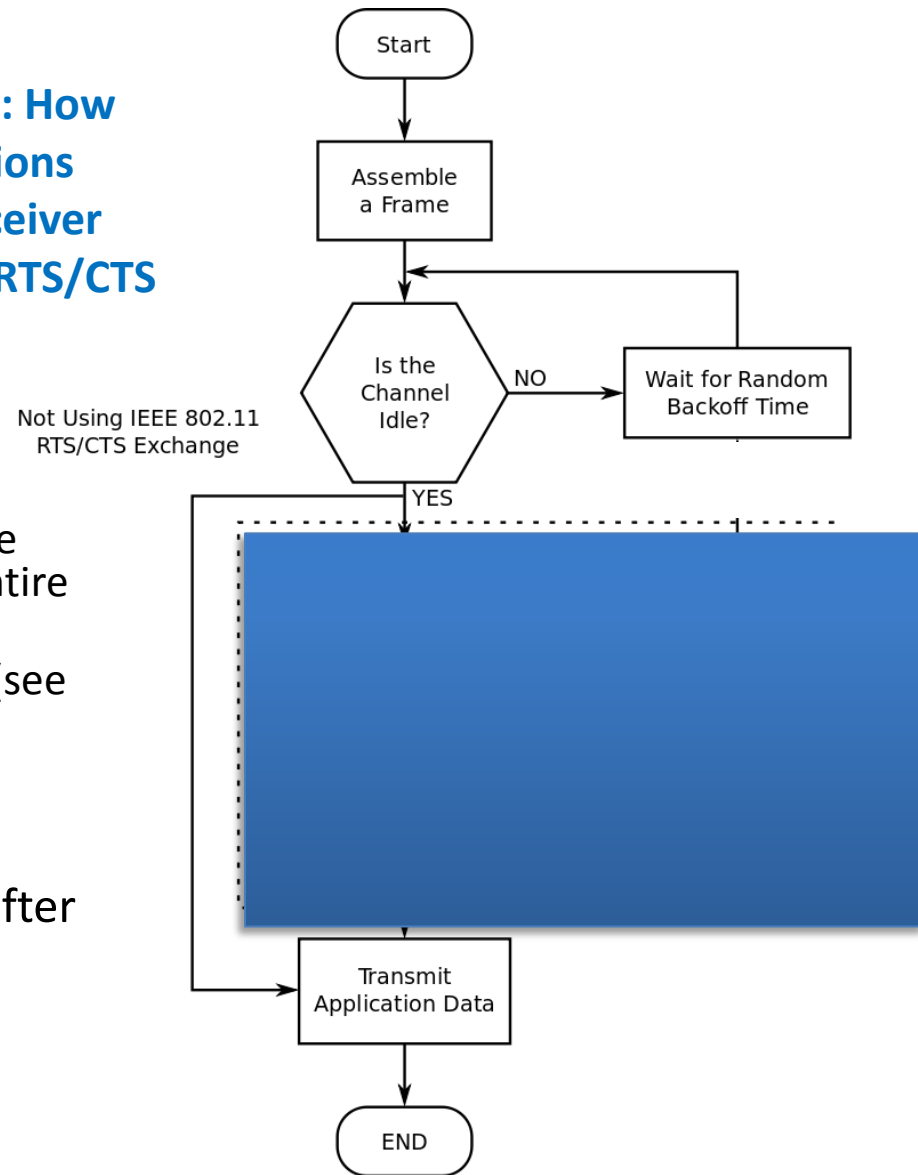


- Need a solution that limits collisions
 - RTS/CTS in CSMA/CA
- Contrary Effect: Exposed Station Problem
 - Imagine B sending to A, C wants to send to D (out of range of B)
 - C waits for A,B finishing their communication (-> unnecessary)

CSMA/CA Collision Avoidance

Q5: Consider the IEEE 802.11 MAC Protocol: How does CSMA/CA tackle the problem of collisions (what steps are taken at the sender and receiver respectively)? What is the idea behind the RTS/CTS concept?

- **Sender:**
 - Sense channel
 - If idle for a certain amount of time (802.11: DIFS, $\sim 50 \mu\text{s}$) transmit entire frame
 - If busy, start **exponential backoff** (see last weeks exercise)
- **Receiver:**
 - If frame received OK, return ACK after waiting a certain amount of time (802.11: SIFS, $\sim 10 \mu\text{s}$)
 - Hidden terminal problem



CSMA/CA RTS/CTS

- Goal: Avoid collisions of **large** data frames
- Idea:
 - Use reservation of channel instead of random access
 - Allow collisions of reservation packets (**small!**)
 - Only reservation packets collide, no data frames!
- Solution: Sender transmits Request-To-Send (RTS) to BS, BS broadcasts Clear-To-Send (CTS) as answer (notifies other nodes in range that channel is busy)

