

# Computer Networks Homework #11

January 28th 2016

# Exercise Exam + Q&A

- Exercise exam
  - Available in wiki
  - Intended for self-study; there will be no answer sheet or exercise session
- Question and Answer Session
  - **February 4<sup>th</sup> 2016**
  - Entirely for your benefit!
  - If there are no questions, there will be no answers
  - If you want a well prepared answer, please send us an email in advance

# NetSec

- What are the security concerns network security is targeting at? What main areas of protection does network security cover?

Confidentiality

Authentication

Message integrity

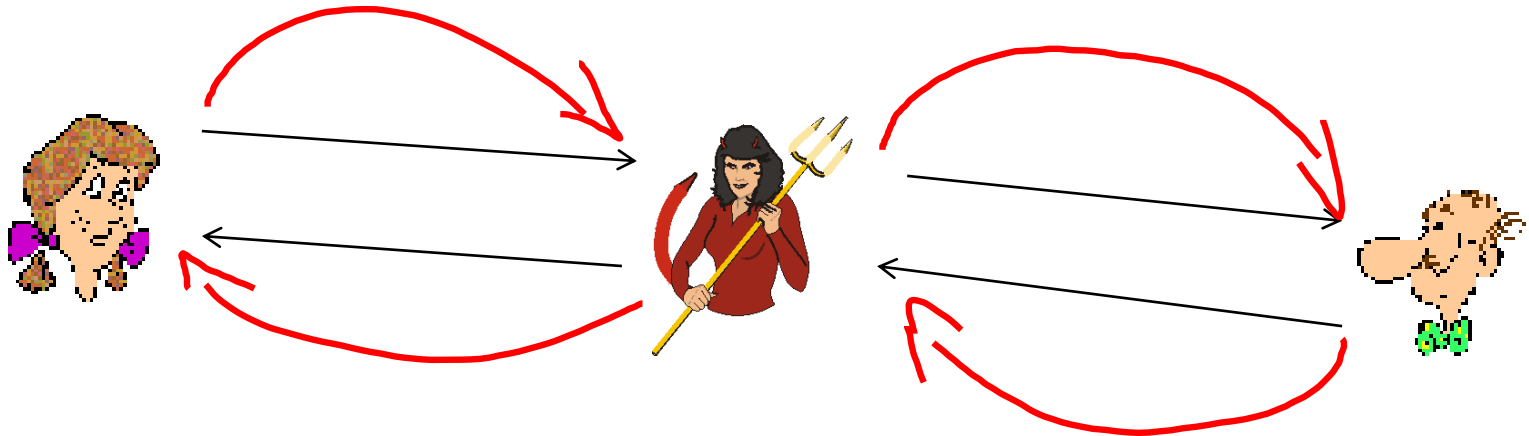
Access and availability

# Cryptography

- What are the two main types of cryptography?
- **Symmetric crypto** (encryption + decryption with the same key): DES, 3DES, AES etc.
- **Asymmetric crypto** (enc and dec with different keys): RSA, Public/Private keying, Diffie-Hellman

# Authentication

- What is a man-in-the-middle attack? Is public key cryptography save against that type of attack?



- Asymmetric keying only helpful if public keys are pre-known or certificate bound.

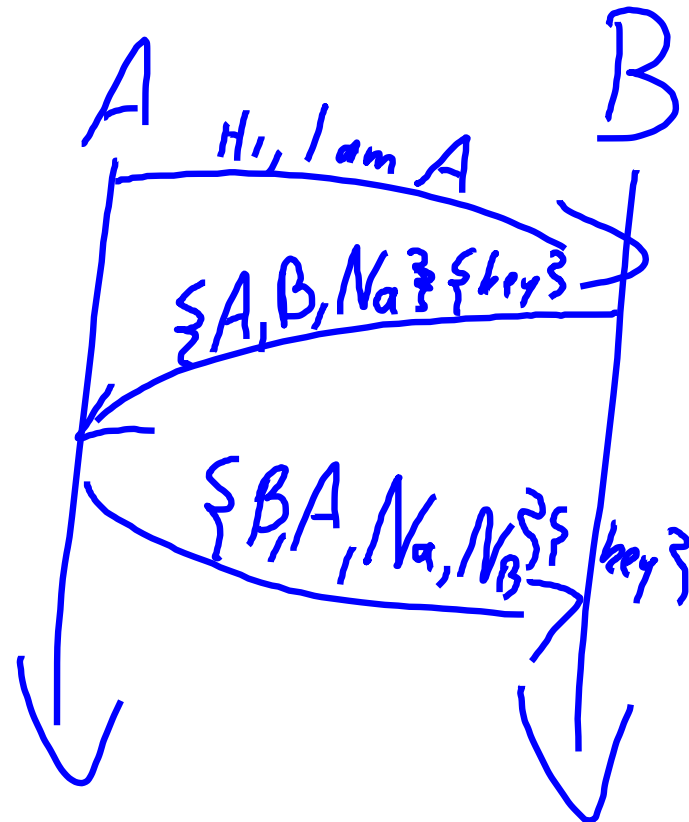
# Authentication

- What other tricks does attackers use to overcome authentication protection? Please explain using the AP protocols presented in the lecture.
- AP 1.0/2.0): Just faking IDs (“I am Alice”) or spoofing an IP address
- Often record and playback attacks as in AP 3.0/3.1

# Nonces

- What is the purpose of a nonce in an end-point authentication protocol?
  - Brings freshness
  - Prevents replay attacks
  - Example:

$\{X\} \xrightarrow{\{key\}}$



# RSA

- Perform an RSA encryption and decryption with  $p=7$  and  $q=11$  with the word “Telematics”.

$n=7*11=77$  (prime factors 7, 11)

$z=(7-1)(11-1)=60$  (prime factors 2, 2, 3, 5)

$e$  needs to be chosen in a way, that it has no common prime factors with  $z$

$e=7$

now we search for a  $d$  with  $e * d - 1 \pmod{z} = 0$ . With  $d=43$  we have

$e*d-1 \pmod{60} = 300 \pmod{60} = 0$



# RSA

$m < n$  (m can be very large!)

$$PK = \{e, n\}$$

$$SK = \{d, n\}$$

*Private*

Klartext		$m^e$	chiffre= $m^e$ mod n		$c^d$ (here: chiffre <sup>46</sup> )	$c^d$ mod n
a	1	1	1		1	1
b	2	128	51			
c	3	2187	31	13444753212776963019174122373997438185440200300120230113873520991		3
d	4	16384	60			
E	5	78125	47	794708560552308362507026214655083140659880205559381016431673633560574223		5
F	6	279936	41			
G	7	823543	28			
H	8	2097152	57			
i	9	4782969	37	27081588506598106040982953896258749653831334409506086433262944331453		9
j	10	10000000	10			
k	11	19487171	11			
l	12	35831808	12	25397652694505813866070015990659936347412758528		12
m	13	62748517	62	118261299920216034323567158324881157722618355000741423528102151243191317168128		13
n	14	105413504	42			
o	15	170859375	71			
p	16	268435456	58			
q	17	410338673	52			
r	18	612220032	39			
s	19	893871739	68	6278895373298528368344913294912019325279912443533041880115104685557599470354432		19
t	20	1280000000	48	1965048198399560713177500537391830916254451560885426333004585474449211392		20
u	21	1801088541	21			
v	22	2494357888	22			
w	23	3404825447	23			
x	24	4586471424	73			
y	25	6103515625	53			
z	26	8031810176	5			

Telematics = 48 47 12 47 62 01 48 37 68

We are encrypting letter by letter, remember cipher algos and consider large m!

# Hashes

- What is the conceptual difference between a crypto-hash function and other hash functions?
  - computationally infeasible to find two different messages,  $x$ ,  $y$  such that  $H(x) = H(y)$ 
    - equivalently: given  $m = H(x)$ , ( $x$  unknown), can not determine  $x$ .
- SHA-1, MD5 operate without a shared secret
- Additionally, key based Hash-based MACs (HMACs) HMAC-MD5 or HMAC-SHA1 available e.g. for signatures

# Thank you

Any questions?