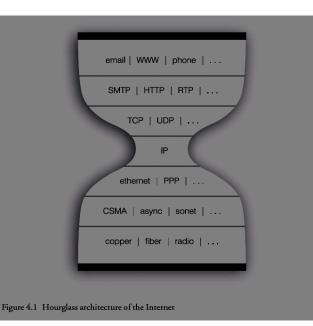
### Exercise 4

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•Why is the Internet Protocol (IP) described as the "narrow waist" of the network stack?

## IP as Narrow Waist

- "IP over anything, anything over IP"
  Single common tie between multiple protocols
- Innovation in upper/lower layers
- •Makes changes to IP difficult (IPv6 ...)



•What are the two key functions of the network layer, that each router performs? Please explain the difference between them.

## Key Functionalities of a Router

#### Network Layer Functions:

Naming

Routing

Abstraction

•Forwarding: move packets from router's input to an appropriate router output

 Routing: determine route taken by packets from source to dest.

•A routing protocol determines the forwarding table

•Assume you have a 3,000 byte long datagram which needs to be fragmented for a 1,400 bytes MTU.

## **IP Datagram Fragmentation**

3000 byte datagram, 1396byte MTU.
One datagram is fragmented into multiple smaller datagrams...

| Datagram No. | Length            | Frag. Flag | Offset =<br>(MTU-Header data<br>/ 8) |
|--------------|-------------------|------------|--------------------------------------|
| 1            | 1396<br>(1376+20) | 1          | 0                                    |
| 2            | 1396<br>(1376+20) | 1          | 172                                  |
| 3            | 248<br>(228+20)   | 0          | 344                                  |
| 4            |                   |            |                                      |

## •Convert the following IP addresses into their binary notion:

.192.168.0.1

## IP Address Conversion (Decimal to Binary)

•TIP: Make yourself a table:

| Power | 2^7 | 2^6 | 2^5 | 2^4 | 2^3 | 2^2 | 2^1 | 2^0 |
|-------|-----|-----|-----|-----|-----|-----|-----|-----|
| Value | 128 | 64  | 32  | 16  | 8   | 4   | 2   | 1   |
| Rest  |     |     |     |     |     |     |     |     |
| Bit   |     |     |     |     |     |     |     |     |

#### ∘For each octet:

Put octet number into first "rest" cell

oBit = (value >= rest ? 1 : 0)

Restnext = Restprev – Bitprev x Valueprev

oRinse and Repeat

## IP Address Conversion (Example)

 $\circ$  First octet of 192.168.0.1

| Power | 2^7 | 2^6 | 2^5 | 2^4 | 2^3 | 2^2 | 2^1 | 2^0 |
|-------|-----|-----|-----|-----|-----|-----|-----|-----|
| Value | 128 | 64  | 32  | 16  | 8   | 4   | 2   | 1   |
| Rest  | 64  | 0   | 0   | 0   | 0   | 2   | 0   | 0   |
| Bit   | 1   | 1   | 0   | 0   | 0   | 0   | 0   | 0   |

## **Converting IP Addresses**

.192.168.0.1

 $\bullet 11000000 \ 10101000 \ 00000000 \ 00000001$ 

## Convert the following IP address into it's decimal notion

 $\bullet 1110001110000110000111110101010$ 

# IP Address Conversion (Binary to Decimal)

•Make yourself a table:

| Power | 2^7 | 2^6 | 2^5 | 2^4 | 2^3 | 2^2 | 2^1 | 2^0 |
|-------|-----|-----|-----|-----|-----|-----|-----|-----|
| Value | 128 | 64  | 32  | 16  | 8   | 4   | 2   | 1   |
| Bit   |     |     |     |     |     |     |     |     |
| Sum   |     |     |     |     |     |     |     |     |

•For each octet:

 $_{\odot}\text{Fill}$  the "Bit" row with the bits of the octet

•Fill the sum row:

Sumnext = Sumprev + Bitprev x Valueprev

# IP Address Conversion (Binary to Decimal)

#### oOctet 11100011:

| Power | 2^7 | 2^6 | 2^5 | 2^4 | 2^3 | 2^2 | 2^1 | 2^0 |
|-------|-----|-----|-----|-----|-----|-----|-----|-----|
| Value | 128 | 64  | 32  | 16  | 8   | 4   | 2   | 1   |
| Bit   | 1   | 1   | 1   | 0   | 0   | 0   | 1   | 1   |
| Sum   | 128 | 192 | 224 | 224 | 224 | 224 | 226 | 227 |

11100011100001100000111110101010227.134.15.170

•A provider has been assigned the network 128.30.0.0/23 and wants to divide it among three customers. Customer A needs to accommodate up to 220 hosts, customer B needs to accommodate up to 110 hosts and customer C needs to accommodate up to 80 hosts. Please fill the following table with the details of the subnetworks that the provider can create to fit its customers' needs.

## **Subnet calculations**

 Subnet calculations are used to break a given network into smaller pieces

A (sub-) network mask shows how many bits
 of an IP address denote the network

oDecimal: /17

•Hexadecimal: 255.255.128.0

## Subnet calculations (Example)

oGiven network: 128.30.0.0/17

•Wanted: Four sub networks

First step: Find new subnet mask

 $_{\circ}$ To address four networks we need at least two bits (2^2 = 4).

• The new subnet mask is 17+2 = 19

Second step: Find new network addresses (see next slide)

 Third step: Calculate data for new networks (see homework)

## Subnet calculations (example)

New netmask: 19 (= 255.255.224.0) 11111111111111111100000.00000000

=> New network 1: 128.30.0.0/19 1000000.00011110.0000000.00000000

=> New network 2: 128.30.32.0/19 1000000.00011110.00100000.00000000

=> New network 3: 128.30.64.0/19
1000000.00011110.01000000.00000000

=> New network 4: 128.30.96.0/19 1000000.00011110.01100000.00000000

Number of hosts:  $2^{13} - 2 = 8,190$ 

## Subnet calculation (homework)

A provider has been assigned the network 128.30.0.0/23 and wants to divide it among three customers. Customer A needs to accommodate up to 220 hosts, customer B needs to accommodate up to 110 hosts and customer C needs to accommodate up to 80 hosts. Please fill the following table with the details of the subnetworks that the provider can create to fit its customers' needs.

| Subnet<br>No. | Network Address | Netmask         | Host Range                     | No. of<br>Hosts |
|---------------|-----------------|-----------------|--------------------------------|-----------------|
| 1<br>Cust. A  | 128.30.0.0/24   | 255.255.255.0   | 128.30.0.1 –<br>128.30.0.254   | 254             |
| 2<br>Cust B   | 128.30.1.0/25   | 255.255.255.128 | 128.30.1.1 –<br>128.30.1.126   | 126             |
| 3<br>Cust C   | 128.30.1.128/25 | 255.255.255.128 | 128.30.1.129 –<br>128.30.1.254 | 126             |
|               |                 |                 |                                |                 |

•Consider IP addresses: How does a host get an IP address? How does a network get the subnet part of an IP address? How does a provider get a block of IP addresses? What is the principle behind these procedures?

## **IP Address Allocation - Host**

•DHCP

Dynamically gets an IP address on joining the network

Allows reuse of addresses (address only reserved while online)

•Protocol: DHCP discover  $\rightarrow$  offer  $\rightarrow$  request  $\rightarrow$  ack

•More details: see lecture slides

## **IP Address Allocation - Network**

- Allocation of a portion of the providers ISP address space
- •e.g., provider net 200.23.16.0/20
- Possible allocated subnet: 200.23.30.0/23

## **IP Address Allocation - Provider**

- ICANN (Internet Corporation for Assigned Names and Numbers)
- •Global allocation of addresses to ISPs
- ISPs then reallocate their addresses to subnets/customers (see previous slides)
- However: Shortage of IPv4 addresses → Most blocks occupied

•What problem is tackled by Network Address Translation (NAT)? Please briefly describe what NAT does.

## Network Address Translation (NAT)

- IPv4: Address shortage
- •NAT: One network (of an arbitrary number of hosts) has only one IP address (NAT enabled router) that is accessible from the internet
- •The remaining hosts are addressed internally
- •Use port numbers to decide which host the datagram is destined to, mapping inside NAT table
- •NAT is often considered a "dirty fix" to the address shortage issue ( $\rightarrow$  IPv6)

•Consider IPv6 What are the main differences between IPv4 and IPv6? What are two approaches towards the transition between IPv4 and IPv6?

## IPv4 vs IPv6 - Differences

•Address space: IPv4 2^32, IPv6 2^128

•IPv6: Fixed header length, additional information needs to be stored in additional headers

•IPv6: No packet fragmentation supported, fragmentation is moved to the sending host

•IPv6: No header checksum, error detection on layer 4/2

••••

## IPv4 to IPv6 - Migration

- There is no "flag day" on which IPv4 routers are replaced by IPv6 routers.
- Not all routers can be upgraded simultaneously
- Rather a slow process of transition
- How to achieve this transition, i.e., a mixed, concurrent operation of IPv4 and IPv6 routers?

## IPv4 and IPv6 together

Two different possibilities

•Tunneling: IPv6 datagram is carried as payload in IPv4 datagram between IPv4 routers; IPv6 routers then decapsulate IPv6 datagram.

•Dual Stack: Routers can do both, IPv4 and IPv6; direct connection between same protocol clients (IPv4  $\rightarrow$  IPv4, IPv6  $\rightarrow$  IPv6); can be used together with tunneling