

Computer Networks Homework #11

January 26th 2012

Announcements

- Final exam: Thursday 09.02.2012
 - 10:00 -12:00 : GZG - MN08
- Language: English + German, answers possible in both languages
- No additional resources (calculator etc.) allowed. Just bring pens ;).
- Register with FlexNow till February 2nd!

Exercise Exam + Q&A

- Exercise exam
 - Available in wiki
 - Intended for self-study; there will be no answer sheet or exercise session
- Question and Answer Session
 - **February 2nd 2012**
 - Entirely for your benefit!
 - If there are no questions, there will be no answers
 - If you want a well prepared answer, please send us an email in advance

NetSec

- What are the security concerns network security is targeting at? What main areas of protection does network security cover?

Confidentiality

Authentication

Message integrity

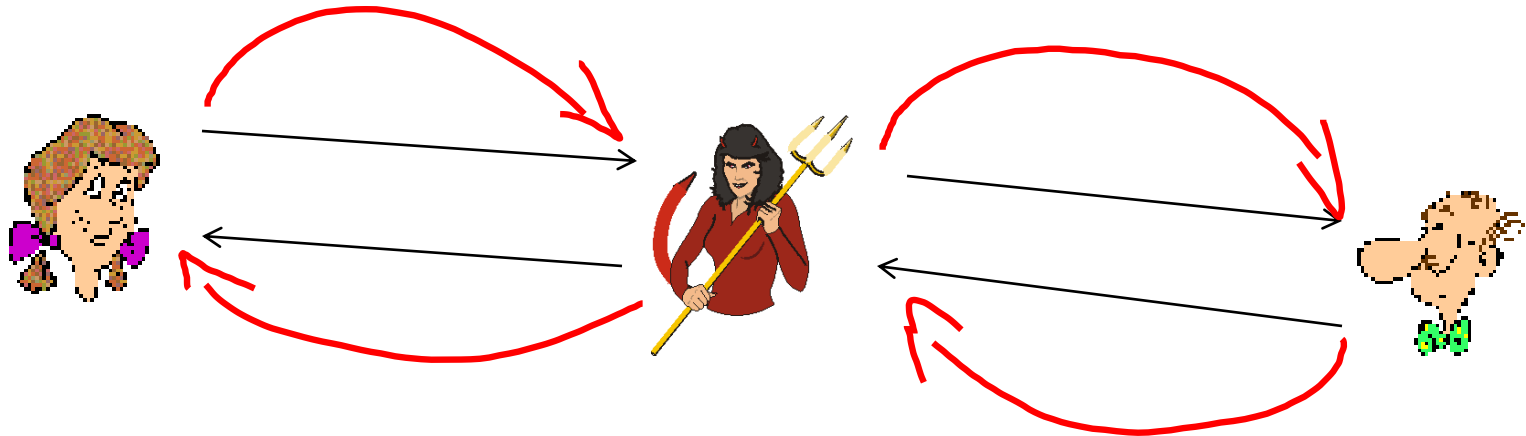
Access and availability

Cryptography

- What are the two main types of cryptography?
- **Symmetric crypto** (encryption + decryption with the same key): DES, 3DES, AES etc.
- **Asymmetric crypto** (enc and dec with different keys): RSA, Public/Private keying, Diffie-Hellman

Authentication

- What is a man-in-the-middle attack? Is public key cryptography save against that type of attack?



- Asymmetric keying only helpful if public keys are pre-known or certificate bound.

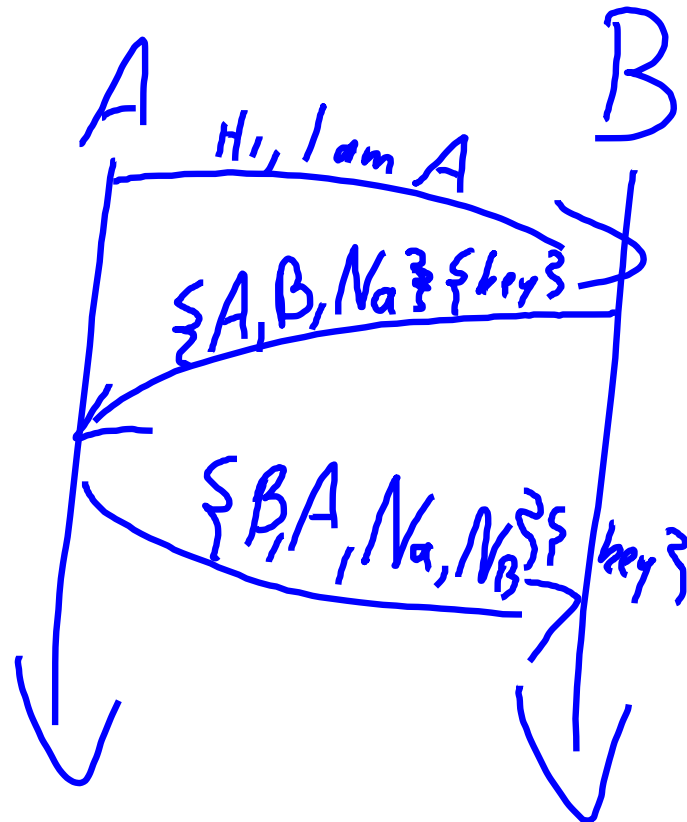
Authentication

- What other tricks does attackers use to overcome authentication protection? Please explain using the AP protocols presented in the lecture.
- AP 1.0/2.0): Just faking IDs (“I am Alice”) or spoofing an IP address
- Often record and playback attacks as in AP 3.0/3.1

Nonces

- What is the purpose of a nonce in an end-point authentication protocol?
 - Brings freshness
 - Prevents replay attacks
 - Example:

$\{X\} \rightarrow \{key\}$



RSA

- Perform an RSA encryption and decryption with $p=7$ and $q=11$ with the word “Telematics”.

$n=7*11=77$ (prime factors 7, 11)

$z=(7-1)(11-1)=60$ (prime factors 2, 2, 3, 5)

e needs to be chosen in a way, that it has no common prime factors with z

$e=7$

now we search for a d with $e * d - 1 \text{ mod } z = 0$. With $d=43$ we have

$e*d-1 \text{ mod } 60 = 300 \text{ mod } 60 = 0$

RSA

$m < n$ (m can be very large!)

$$PK = \{e, n\}$$

$$SK = \{d, n\}$$

Private

Klartext		m^e	chiffre= m^e mod n		c^d (here: chiffre 46)	c^d mod n
a	1	1	1		1	1
b	2	128	51			
c	3	2187	31	13444753212776963019174122373997438185440200300120230113873520991		3
d	4	16384	60			
E	5	78125	47	794708560552308362507026214655083140659880205559381016431673633560574223		5
F	6	279936	41			
G	7	823543	28			
H	8	2097152	57			
i	9	4782969	37	27081588506598106040982953896258749653831334409506086433262944331453		9
j	10	10000000	10			
k	11	19487171	11			
l	12	35831808	12	25397652694505813866070015990659936347412758528		12
m	13	62748517	62	118261299920216034323567158324881157722618355000741423528102151243191317168128		13
n	14	105413504	42			
o	15	170859375	71			
p	16	268435456	58			
q	17	410338673	52			
r	18	612220032	39			
s	19	893871739	68	6278895373298528368344913294912019325279912443533041880115104685557599470354432		19
t	20	1280000000	48	1965048198399560713177500537391830916254451560885426333004585474449211392		20
u	21	1801088541	21			
v	22	2494357888	22			
w	23	3404825447	23			
x	24	4586471424	73			
y	25	6103515625	53			
z	26	8031810176	5			

Telematics = 48 47 12 47 62 01 48 37 68

We are encrypting letter by letter, remember cipher algos and consider large m!



Hashes

- What is the conceptual difference between a crypto-hash function and other hash functions?
 - computationally infeasible to find two different messages, x , y such that $H(x) = H(y)$
 - equivalently: given $m = H(x)$, (x unknown), can not determine x .
- SHA-1, MD5 operate without a shared secret
- Additionally, key based Hash-based MACs (HMACs) HMAC-MD5 or HMAC-SHA1 available e.g. for signatures

Thank you

Any questions?