# Telematics Homework \#4 

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## What kind of switching fabrics are there and how do they work?

- Bus switching
- Fabric connects ports by means of a bus that is shared among all ports
- Memory switching
- Fabric connects ports by means of a common memory that is used by all ports
- Crossbar switching
- Fabric can connect any input port to any output port directly


## When and where does buffering occur? What are its effects?

- Input ports:
- Fabric slower than incoming traffic
- Output ports:
- Datarate from fabric is faster than outgoing data rate
- Buffering introduces queueing delays and ultimately leads to loss


## When and where does head-ofline blocking occur?

- When a packet is queued at an input port, following packets must wait until that packet is processed


# In an IP datagram: what is the header checksum for and where is it calculated? 

- Used for error-checking of the header
- By intermediate routers
- By the destination
- Calculated
- At the source
- At every intermediate router after decrementing the time-to-live value


## Where does IP fragmentation occur?

- IP fragmentation occurs in intermediate routers when a datagram is bigger than the maximum transmission unit (MTU) of the following network segment


## Fragmentation

- Assume you have a 2,000 byte long datagram which needs to be fragmented for a 1,000 bytes MTU. Please fill the following table with the data of the resulting datagrams.

| Datagram No. | Length | Frag. Flag | Offset |
| ---: | :--- | :--- | :--- |
| 1 | 996 | 1 | 0 |
| 2 | $(976+20)$ |  | 122 |
| 3 | 996 | 1 | 244 |
| 4 | 48 |  | 0 |

## Tip: IP Address Conversion (Decimal to Binary)

- Make yourself a table:

| Power | $\mathbf{2}^{7}$ | $\mathbf{2}^{6}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{4}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{2}$ | $\mathbf{2}^{1}$ | $\mathbf{2}^{0}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Rest |  |  |  |  |  |  |  |  |
| Bit |  |  |  |  |  |  |  |  |

- For each octet:
- Put octet number into first „rest" cell
- Bit = (value >= rest? $1: 0$ )
- Rest $_{\text {next }}=$ Rest $_{\text {prev }}-$ Bit $_{\text {prev }} \times$ Value $_{\text {prev }}$
- Rinse and Repeat


## Tip: IP Address Conversion (Example)

- First octet of 66.135.207.138:

| Power | $\mathbf{2}^{7}$ | $\mathbf{2}^{\mathbf{6}}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{4}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{\mathbf{2}}$ | $\mathbf{2}^{\mathbf{1}}$ | $\mathbf{2}^{0}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Rest |  |  |  |  |  |  |  |  |
| Bit |  |  |  |  |  |  |  |  |

- Result:

| Power | $\mathbf{2}^{\mathbf{7}}$ | $\mathbf{2}^{\mathbf{6}}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{\mathbf{4}}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{\mathbf{2}}$ | $\mathbf{2}^{\mathbf{1}}$ | $\mathbf{2}^{\mathbf{0}}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Rest | 66 | 66 | 2 | 2 | 2 | 2 | 2 | 0 |
| Bit | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 |

# Convert the following IP addresses into their binary notion 

- 66.135.207.138
- 01000010.10000111.11001111.10001010
- 192.35.225.7
- 11000000.00100011.11100001.00000111


## Tip: IP Address Conversion (Binary to Decimal)

- Make yourself a table:

| Power | $\mathbf{2}^{7}$ | $\mathbf{2}^{6}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{4}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{2}$ | $\mathbf{2}^{1}$ | $\mathbf{2}^{0}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Bit |  |  |  |  |  |  |  |  |
| Sum |  |  |  |  |  |  |  |  |

- For each octet:
- Fill the „Bit" row with the bits of the octet
- Fill the sum row:

Sum $_{\text {next }}=$ Sum $_{\text {prev }}+$ Bit $_{\text {prev }} \times$ Value $_{\text {prev }}$

## Tip: IP Address Conversion (Example)

- Octet 10000110:

| Power | $\mathbf{2}^{\mathbf{7}}$ | $\mathbf{2}^{\mathbf{6}}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{\mathbf{4}}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{\mathbf{2}}$ | $\mathbf{2}^{\mathbf{1}}$ | $\mathbf{2}^{\mathbf{0}}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Bit |  |  |  |  |  |  |  |  |
| Sum |  |  |  |  |  |  |  |  |

- Result:

| Power | $\mathbf{2}^{\mathbf{7}}$ | $\mathbf{2}^{\mathbf{6}}$ | $\mathbf{2}^{\mathbf{5}}$ | $\mathbf{2}^{\mathbf{4}}$ | $\mathbf{2}^{\mathbf{3}}$ | $\mathbf{2}^{\mathbf{2}}$ | $\mathbf{2}^{\mathbf{2}}$ | $\mathbf{2}^{\mathbf{0}}$ |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| Value | 128 | 64 | 32 | 16 | 8 | 4 | 2 | 1 |
| Bit | 1 | 0 | 0 | 0 | 0 | 1 | 1 | 0 |
| Sum | 128 | 128 | 128 | 128 | 128 | 132 | 134 | $\mathbf{1 3 4}$ |

## Convert the following IP addresses into their decimal notion

- 10000110.01001100.01010001.00011001
- 134.76.81.25
-11011000.10011110.01010111.00010111
- 216.158.87.23


## Tip: Subnet calculations

- Subnet calculations are used to break a given network into smaller pieces
- A (sub-) network mask shows how many bits of an IP address denote the network
- Decimal: /17
- Binary: 11111111.11111111.10000000.00000000
- Hexadecimal: 255.255.128.0


## Tip: Subnet calculations

- Given address: 128.30.10.0
- 10000000.00011110.00001010.00000000
- Given netmask: 17 (= 255.255.128.0)

。11111111.11111111.10000000.00000000

- => Network: 128.30.0.0/17
- 10000000.00011110.00000000.00000000
- => Broadcast: 128.30.127.255
- 10000000.00011110.01111111.11111111
- => First host: 128.30.0.1
- 10000000.00011110.00000000.00000001
- => Last host: 128.30.127.254
- 10000000.00011110.01111111.11111110
- Number of hosts: $2^{15}-2=32,766$


## Tip: Subnet calculations (Example)

- Given network: 128.30.0.0/17
- Wanted: Four sub networks
- First step: Find new subnet mask
- To address four networks we need at least two bits $\left(2^{2}=4\right)$.
- The new subnet mask is $17+2=19$
- Second step: Find new network addresses (see next slide)
- Third step: Calculate data for new networks (see previous slide)


# Tip: Subnet calculations (Example) 

- New netmask: 19 (= 255.255.224.0)
- 11111111.11111111.11100000.00000000
- => New network 1: 128.30.0.0/19
-10000000.00011110.00000000.00000000
- => New network 2: 128.30.32.0/19
- 10000000.00011110.00100000.00000000
- => New network 3: 128.30.64.0/19
- 10000000.00011110.01000000.00000000
- => New network 4: 128.30.96.0/19
- 10000000.00011110.01100000.00000000
- Number of hosts: $2^{13}-2=8,190$


## Subnet calculation

- A provider has been assigned the network 128.30.0.0/17 and wants to divide it into four subnets. Please fill the following table with the data of the resulting sub networks.

| Subnet <br> No. | Network <br> Address | Netmask | Host Range | No. of <br> Hosts |
| ---: | ---: | ---: | ---: | ---: |
| 1 | $128.30 .0 .0 / 19$ | 255.255 .224 .0 | $128.30 .0 .1-$ <br> 128.30 .31 .254 | 8,190 |
| 2 | $128.30 .32 .0 / 19$ | 255.255 .224 .0 | $128.30 .32 .1-$ <br> 128.30 .64 .254 | 8,190 |
| 3 | $128.30 .64 .0 / 19$ | 255.255 .224 .0 | $128.30 .64 .1-$ <br> 128.30 .95 .254 | 8,190 |
| 4 | $128.30 .96 .0 / 19$ | 255.255 .224 .0 | $128.30 .96 .1-$ | 8,190 |

## What is the "NAT traversal problem" and how can it be solved?

- NAT traversal problem:
- A server behind a NAT cannot be reached directly because no mapping between external and internal address exists
- Solution 1: Statically configured mapping
- Solution 2: Dynamically configured mapping (e.g. by using UPnP)


## What is the "NAT traversal problem" and how can it be solved? (cont'd)

- Solution 3: Relaying
- Node behind NAT establishes connection to relay
- External node connects to relay
- Relay forwards packets between nodes


# What are the main differences between IPv4 and IPv6? 

- Bigger address space in IPv6 (128 bit vs. 32 bit)
- Fixed-length 40 byte header in IPv6
- No fragmentation allowed in IPv6
- No header checksum in IPv6
- Options outside of header in IPv6
- New version of ICMP in IPv6 (ICMPv6)


## How can an IPv4 host communicate with an IPv6 host?

- Trick question:
- An IPv4 host cannot directly communicate with an IPv6 host
- Tunneling:
- Allows to connect IPv6 hosts with IPv4 hosts in the middle
- Dual stack:
- Use IPv4 to communicate with IPv4 hosts and IPv6 to communicate with IPv6 hosts
- There will be no trick questions in the written examination!


## Thank you

Any questions?

