# Telematics Homework \#11 

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## Announcements

o Final exam: Thursday 04.02.2010 - 10:00-12:00 : GZG - MN08
o Language: English + German, answers possible in both languages
o No additional resources (calculator etc.) allowed. Just bring pens ;).

## Exercise Exam + Q\&A

o Exercise exam

- Available in wiki and StudIP
- Intended for self-study; there will be no answer sheet or exercise session
- Give us feedback on size and difficulty!
- Question and Answer Session
o 28 January 2010, 10:15h (until all questions are answered)
- Entirely for your benefit!
- If there are no questions, there will be no answers
o If you want a well prepared answer, please send us an email in advance


## NetSec

o What are the security concerns network security is targeting at? What main areas of protection does network security cover?

Confidentiality<br>Authentication<br>Message integrity<br>Access and availability

## Cryptopgraphy

o What are the two main types of cryptography?
o Symmetric crypto (encryption + decryption with the same key): DES, 3DES, AES etc.
o Asymmetric crypto (enc and dec with different keys): RSA, Public/Private keying, DiffieHellman

## Authentication

o What is a man-in-the-middle attack? Is public key cryptography save against that type of attack?

o Asymmetric keying only helpful if public keys are pre-known or certificate bound.

## Authentication

o What other tricks does attackers use to overcome authentication protection? Please explain using the AP protocols presented in the lecture.
o AP 1.0/2.0): Just faking IDs ("I am Alice") or spoofing an IP address
o Often record and playback attacks as in AP 3.0/3.1

Nonces
Tunnel

- What is the purpose of a nonce in an and point authentication protocol?
- Brings freshness
- Prevents replay attacks
- Example:

$$
\{X\}\{\text { key }\}
$$



## RSA

- Perform an RSA encryption and decryption with $\mathrm{p}=\boldsymbol{3}$ and $q=\$ 1$ with the word "Telematics".
$\mathrm{n}=7 * 11=77$ (prime factors 7,11 )
$z=(7-1)(11-1)=60$ (prime factors $2,2,3,5)$
e needs to be chosen in a way, that it has no common prime factors with $z$ $e=7$
now we search for a $d$ with $e^{*} d-1 \bmod z=0$. With $d=43$ we have $e^{\star} d-1 \bmod 60=300 \bmod 60=5$

$$
P K=\{e, n\}
$$



| Klartext |  | $\mathrm{m}^{\wedge} \mathrm{e}$ | chiffre $=\mathrm{m}^{\wedge} \mathrm{e}$ $\bmod n$ | $c^{\wedge} \mathrm{d}$ | $\mathrm{c}^{\wedge} \mathrm{d} \mathrm{mod} \mathrm{n}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| a | 1 | 1 | 1 | 1 | 1 |
| b | 2 | 128 | 51 |  |  |
| c | 3 | 2187 | 31 | 13444753212776963019174122373997438185440200300120230113873520991 | 3 |
| d | 4 | 16384 | 60 |  |  |
| E |  | 78125 | 47 | 794708560552308362507026214655083140659880205559381016431673633560574223 | 5 |
| F |  | 279936 | 41 |  | - |
| G | 7 | 823543 | 28 |  |  |
| H | 8 | 2097152 | 57 |  |  |
| i | 9 | 4782969 | 37 | 27081588506598106040982953896258749653831334409506086433262944331453 | 9 |
| j | 10 | 10000000 | 10 |  |  |
| k | 11 | 19487171 | 11 |  |  |
| 1 | 12 | 35831808 | 12 | 25397652694505813866070015990659936347412758528 | 12 |
| m | 13 | 62748517 | 62 | 118261299920216034323567158324881157722618355000741423528102151243191317168128 | 13 |
| n | 14 | 105413504 | 42 |  |  |
| 0 | 15 | 170859375 | 71 |  |  |
| p | 16 | 268435456 | 58 |  |  |
| q | 17 | 410338673 | 52 |  |  |
| r | 18 | 612220032 | 39 |  |  |
| S | 19 | 893871739 | 68 | 6278895373298528368344913294912019325279912443533041880115104685557599470354432 | 19 |
| t | 20 | 1280000000 | 48 | 1965048198399560713177500537391830916254451560885426333004585474449211392 | 20 |
| u | 21 | 1801088541 | 21 |  |  |
| v | 22 | 2494357888 | 22 |  |  |
| w | 23 | 3404825447 | 23 |  |  |
| x | 24 | 4586471424 | 73 |  |  |
| y | 25 | 6103515625 | 53 |  |  |
| Z | 26 | 8031810176 | 5 |  |  |

Telematics $=484712476201483768$

[^0]We are encrypting letter by letter, remember cipher algos and consider large m !

## Hashes

- What is the conceptual difference between a crypto-hash function and other hash functions?
o computationally infeasible to find two different messages, x , $y$ such that $H(x)=H(y)$
- equivalently: given $m=H(x)$, ( $x$ unknown), can not determine $x$.
- SHA-1, MD5 operate without a shared secret
o Additionally, key based Hash-based MACs (HMACs) HMAC-MD5 or HMAC-SHA1 available e.g. for signatures


## Thank you

## Any questions?


[^0]:    $\left.\mathbf{N}^{-2} \mathbf{E}^{-1} \mathbf{T}^{-}()\right)$
    WOR-KS

